# WORKING PAPER

OUTLINE OF A DATA MODIFICATIONS SYSTEM FOR AGGREGATION OF COMMODITY FLOWS

Alexander Sarkisov Börje Johansson

November, 1983 WP-83-105



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FOREWORD

This paper has been motivated by the need to develop software and methods for analyzing, modifying and restructuring data on world trade in the Forest Sector Project. The paper outlines a system of approaches to (i) modifying, correcting, and aggregating initial data on bilateral trade flows and (ii) disaggregating aggregate flows obtained from model projections. The approaches suggested have their origin in experiences with multiregional modeling. The computer programs corresponding to the methods described in the paper have been, and are being, further developed by Alexander Sarkisov who participated in the Regional and Urban Development Group within the YSSP program.

> BÖrje Johansson Acting Leader Regional & Urban Development Group

October, 1983

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# OUTLINE OF A DATA MODIFICATIONS SYSTEM FOR AGGREGATION OF COMMODITY FLOWS

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#### 1. INTRODUCTION

# 1.1 Trade Flows, Differentiated Competition and Aggregation

This paper is concerned with problems related to the analysis of trade matrices describing bilateral trade flows and associated price and transportation costs structures. Models for generating scenarios of trade patterns have to be formulated in terms of aggregate trade matrices in order to be computationally tractable. Such aggregation is accomplished by forming groups of commodities (commodity groups) and groups of countries (regions). The aggregate structure should ideally have groups which are as homogenous as possible. This is important for three reasons: (i) to make the model more reliable, (ii) to make it possible to interpret model results in a reasonable way, and (iii) eventually to disaggregate the model results into a finer commodity classification and regional partitioning.

This problem of aggregation-disaggregation is a general problem in trade analysis, since all commodity classifications refer to non-homogenous products and all regions are nonhomogenous spatial aggregates. Moreover, the analysis of trade flows brings product differentiation and differentiated competition into focus. According to the classical analysis of Robinson

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(1933) and Chamberlin (1933) every bilateral flow should be looked upon as a differentiated commodity, which implies that prices potentially should be considered as link-specific. Another related issue concerns product development which may imply that a commodity group may contain components which change quality and attributes while others remain unchanged over time.

The kind of problems mentioned above are discussed in this paper both from the viewpoint of methods for (i) disaggregating model results and (ii) analyzing historical data in order to create meaningful aggregation schemes.

#### 1.2 Data Analysis and Data Modification Systems

The earlier introduction motivates the need for a data processing system which can be used to modify and restructure a data base. Moreover, such a data base is often created from many different sources which means that the data analysis and data processing must solve problems of consistency. In particular, the reshaping of data should be problem-specific in order to be purposeful. Since the same data base may be used for several different problems, there is a need for a flexible data processing system which can transform the original data to a problem-relevant form (compare Johansson & Marksjö, 1983). This is described in Figure 1.



Figure 1. Scheme for data transformation.

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The analysis in the paper is based on experiences with the Data Modification System which is currently being developed in the Institute for Systems Studies in Moscow (USSR), (Britkov, V., Sarkisov, A., 1983). Such a system may be described as an interactive artificial intelligence system which is based on an ambition to use a language which is close to the natural language related to the given or specified problem. The main operations of such a system are: (i) consistency checking, (ii) classification, (iii) solving for missing information, (iv) aggregation and disaggregation, etc. Those operations may be summarized by the notions data processing and preliminary data analysis.

The presentation also refers to recent results as regards aggregation analysis (Batten, 1982; Roy, Batten & Lesse, 1982; Lesse, Batty & Batten, 1983). Several of the procedures proposed rely on the existence of algorithms of the type described in Eriksson (1981), Andersson & Persson (1982), and Batten (1982). Moreover, the approaches described are related to Swedish experiences of creating flexible multiregional information systems (Johansson & Marksjö, 1983).

The problem analyzed in the paper is illustrated in Figure 2, where  $x_i^{rs}$  denotes delivery of commodity i from country r to country s, and where  $x_I^{RS}$  has a similar interpretation for aggregates I $\ni$ i, R $\ni$ r and S $\ni$ s



Figure 2. Principal procedure for carrying out aggregationmodeling-disaggregation.

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2.1 Specification of Commodity Flows and Composition of Aggregates

We are studying trade flows of commodities between regions. Our analysis refers in particular to models of world trade, in which a region is a group of countries and a commodity is a group of non-identical traded products. A commodity flow is specified in three basic dimensions by:

- (a) a commodity index, i, signifying to which product class the flow refers
- (b) an origin index, r, signifying from which region the flow is delivered (exported)
- (c) a destination index, s, signifying
   to which region the flow is delivered
   (i.e., into which region it is imported)

As a point of reference we shall introduce a specification of commodities which we call the most disaggregated. This specification consists of (i) the index set for commodities  $i \in II$ , (ii) exporting regions  $r \in \mathbb{R}$ , and (iii) importing regions  $s \in \mathbb{R}$ . Subsets of these indices are denoted by  $I \subset II$ ,  $R \subset \mathbb{R}$ , and  $S \subset \mathbb{R}$ . We shall use I, R, and S to denote sets as well as index numbers. In the first case we can write  $I \in \{I_1, \ldots, I_n\}$  where  $I_i \subset II$ ;  $R \in \{R_1, \ldots, R_m\}$  where  $R_i \subset \mathbb{R}$ ; and  $S \in \{S_1, \ldots, S_k\}$  where  $S_i \subset \mathbb{R}$ . In the second case we have  $I = 1, \ldots, n$ ;  $R = 1, \ldots, m$ ; and  $S = 1, \ldots, k$ . With this we define

> An aggregation scheme I=1,...,n, r=1,...,m, and s=1,...,k is complete if simultaneously  $\cup \{I_j : j=1,...,n\} = II, \cap \{I_j : j=1,...,n\} = \emptyset$  $\cup \{R_j : j=1,...,m\} = IR, \cap \{R_j : j=1,...,n\} = \emptyset$  $\cup \{S_j : j=1,...,k\} = IR, \cap \{S_j : j=1,...,n\} = \emptyset$  (2.1)

Trade flows between origin, r or R, and destination, s or S, can be measured in quantity terms, signified by x, value terms at the origin, signified by V, value terms at the destination, signified by W. Using the indexation in (2.1), a trade flow

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can be characterized by  $(x_i^{rs}, v_i^{rs}, W_i^{rs})$  or  $(x_I^{RS}, v_I^{RS}, W_I^{RS})$ . Let  $c_i^{rs}$  and  $c_I^{RS}$  denote unit transportation costs for the links (r,s) and (R,S). Then we can derive an upper value for transportation costs as follows:

$$c_{i}^{rs} x_{i}^{rs} \leq W_{i}^{rs} - V_{i}^{rs}$$

$$c_{I}^{Rs} x_{I}^{Rs} \leq W_{I}^{Rs} - V_{I}^{Rs}$$
(2.2)

Consider now that we have a model with which we are generating projections (scenarios, forecasts, etc.) for aggregate trade flows (I,R,S). We want these flows to represent the individual disaggregate flows which cannot or should not be studied directly. The reason for this may be that the size of a disaggregated model is intractable, and/or that invariances as regards the behavior of the system are more significant in an aggregate form. Our desire is to construct aggregation schemes that make it possible to disaggregate the projections which we obtain from our model. We also want to know which additional information we need to perform the disaggregation. Moreover, even if no disaggregation is carried through, we need conceptual models to interpret aggregate projections that are due to the aggregation.

We shall not attempt to provide any definite solutions to the problems mentioned above. Instead we shall examine the possibilities of utilizing a data modification system and associated methods for analyzing data and constructing aggregation schemes which are designed for each specific problem.

With regard to spatial aggregation we may identify the following forms:

$$(\mathbf{x}_{i}^{\mathrm{RS}}, \mathbf{v}_{i}^{\mathrm{RS}}, \mathbf{w}_{i}^{\mathrm{RS}}) = \sum_{\substack{r \in \mathbb{R} \\ s \in \mathbb{S}}} (\mathbf{x}_{i}^{\mathrm{rs}}, \mathbf{v}_{i}^{\mathrm{rs}}, \mathbf{w}_{i}^{\mathrm{rs}})$$
(2.3)  
$$(\mathbf{x}_{i}^{\mathrm{R*}}, \mathbf{w}_{i}^{\mathrm{R*}}, \mathbf{w}_{i}^{\mathrm{R*}}) = \sum_{\substack{r \in \mathbb{N} \\ s \in \mathbb{S}}} (\mathbf{x}_{i}^{\mathrm{rs}}, \mathbf{w}_{i}^{\mathrm{rs}})$$
(2.3)

$$(\mathbf{x}_{i}^{R^{*}}, \mathbf{v}_{i}^{R^{*}}, \mathbf{w}_{i}^{R^{*}}) = \sum_{\substack{\boldsymbol{\Gamma} \in \mathbf{R} \\ \mathbf{s} \in \mathbb{R}}} (\mathbf{x}_{i}^{rs}, \mathbf{v}_{i}^{rs}, \mathbf{w}_{i}^{rs})$$
(2.4)

$$(\mathbf{x}_{i}^{*S}, \mathbf{V}_{i}^{*S}, \mathbf{W}_{i}^{*S}) = \sum_{\substack{s \in S \\ r \in \mathbb{R}}} (\mathbf{x}_{i}^{rs}, \mathbf{V}_{i}^{rs}, \mathbf{W}_{i}^{rs})$$
(2.5)

where  $x_i^{R*}$  and  $x_i^{*S}$  represent total supply in region R, and total demand in region S. Commodity aggregation takes the form

$$(\mathbf{x}_{\mathbf{I}}^{\mathrm{rs}}, \mathbf{V}_{\mathbf{I}}^{\mathrm{rs}}, \mathbf{W}_{\mathbf{I}}^{\mathrm{rs}}) = \sum_{i \in \mathbf{I}} (\mathbf{x}_{i}^{\mathrm{rs}}, \mathbf{V}_{i}^{\mathrm{rs}}, \mathbf{W}_{i}^{\mathrm{rs}})$$
(2.6)

The remaining forms of aggregation are obtained by combining (2.6) with (2.3)-(2.5).

# 2.2 Composition of Regional Commodity Groups

The most disaggregate form of regional commodities as given by (2.4) and (2.5) is  $x_i^{r*} = \sum_s x_i^{rs}$  and  $x_i^{*s} = \sum_r x_i^{rs}$ , where the first denotes regional supply (at origin r) and the second regional demand (at destination s). We should observe that in practice  $x_i^{r*}$  is an aggregate summarizing the supply of different firms whose products have non-identical attributes; they may differ with regard to technical characteristics, but also in terms of logistics, guarantee of quality provided by a good name, quickness of service, length of credit, advertisement, etc., (compare Robinson, 1933). These aspects become even more evident when we are forming commodity aggregates

> $\mathbf{x}_{I}^{r*} = \sum_{i \in I} \mathbf{x}_{i}^{r*}$ (2.7)  $\mathbf{x}_{I}^{*S} = \sum_{i \in T} \mathbf{x}_{i}^{*S}$ (2.8)

Using the same type of summation as in (2.7) and (2.8) we can form the price variables  $v_{I}^{k*} = V_{I}^{k*}/x_{I}^{r*}$  and  $w_{I}^{*k} = W_{I}^{*k}/x_{I}^{*k}$ . If there is no price dispersion and competition is perfect the regional price,  $p_{I}^{r}$ , satisfies the "ideal" condition <sup>1</sup>

$$p_{I}^{k} = v_{I}^{k*} = w_{I}^{*k}$$
 (2.9)

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<sup>1)</sup> In order to avoid a jungle of notations we only consider (throughout the paper) prices which are obtained after the effects of tariffs and the like have been removed.

In order to draw conclusions about regional prices of individual products,  $p_i^r$ , from observations of  $p_I^r$ , we must use some assumption about the aggregate commodity I. One may, e.g., assume that I is a composite commodity, which has the following meaning

> Let t denote time. Then a commodity group I is a *composite commodity* if for every  $i \in I$  there exists a positive time-invariant scalar  $\alpha_j^r$  such that we have  $p_j^r(t) = \alpha_j^r p_I^r(t)$ . (2.10)

Suppose that products are developed over time so that the attributes of commodity  $j \in I$  are changed in such a way that the quality is improved. This might give rise to a change process in which  $\alpha_j^r(t)$  increases over time. For another commodity  $i \in I$  we may simultaneously observe that  $\alpha_i^r(t)$  gradually decreases over time. From this we may define

The commodity group I is a composite commodity in a weak sense if for each  $i \in I$  there exist known (2.11) functions  $\alpha_i^r(t) > 0$  such that we have  $p_i^r(t) = \alpha_i^r(t)$  $p_T^r(t)$ .

Suppose that we have information about the two aggregates  $v_I^{r*}$  and  $x_I^{r*}$  and want to decompose this information into the disaggregated elements  $v_i^{r*}$  and  $x_i^{r*}$ , where  $v_I^{r*}$  may be regarded as a proxy for  $p_I^r$ . In order to carry through the decomposition we need the following strong condition

A commodity group I is regionally decomposable if for each  $i \in I$  there exists a pair of known functions  $(\alpha_i^r(t), \beta_i^r(t))$  such that (i)  $p_i^r(t) = \alpha_i^r(t) p_I^r(t)$ , (ii)  $x_i^{r*}(t) = \beta_i^r(t) x_I^{r*}(t)$ , (iii)  $\sum_{i=1}^{r} (t) \beta_i^r(t) = 1$  and  $\sum_{i=1}^{r} (t) = 1$ .

where condition (iii) is a result of the simultaneous fulfilment of (i) and (ii).

Let  $V_{I}^{r*}(t)$  and  $x_{I}^{r*}(t)$  be forecasted values. It means that the regional price of commodity I is known,  $v_{I}^{r*}(t) = V_{I}^{r*}(t)/x_{I}^{r*}(t)$ , and can serve as a proxy for  $p_{I}^{r}(t)$ . Consider then the case when we assume I to be a composite commodity in the weak sense  $(\alpha_{i}^{r}(t)$ are known functions of time). If for this case we also have information (or introduce assumptions) about the maximum supply capacity for individual products  $X_{i}^{r}(t)$  then the fact that values  $\beta_{i}^{r}(t) \geq 0$  and  $\sum_{i} \beta_{i}^{r} = 1$  and thus can be considered as probability distributions, allows us to use the information theory approach to calculate the most probable values of  $\beta_{i}^{r}(t) = \hat{\beta}_{i}^{r}(t)$  by solving for any fixed time point the following minimization problem:<sup>1)</sup>

$$\min \sum_{i} \sum_{i}^{r} \log \beta_{i}^{r}$$
s.t.
$$\beta_{i}^{r} x_{I}^{r} \leq X_{i}^{r}$$

$$\sum_{i} \alpha_{i}^{r} \beta_{i}^{r} = 1$$

$$(2.13)$$

Observe that in this case we know cost functions of individual products  $p_i^r(t) = \alpha_i^r(t) p_I^r(t)$ .

Now for the same forecasted values consider the case when the sectoral change paths  $\beta_i^r(t)$  and additional information about maximum price values for individual products  $p_i^r(max)$  are given. If we substitute new variables  $\xi_i^r(t)$  for  $\alpha_i^{rs}$  so that  $\xi_i^r(t) = \alpha_i^r(t) \beta_i^r(t)$  then the condition for regional decomposition of commodity I can be rewritten:

<sup>1)</sup> The solution to (2.13) could also be called "maximally non-committed with regard to missing information" (Jaynes, 1957).

(i) 
$$p_{i}^{r}(t) = (\xi_{i}^{r}(t)/\beta_{i}^{r}(t)) p_{I}^{r}(t)$$
  
(ii)  $x_{i}^{r*}(t) = \beta_{i}^{r}(t) x_{I}^{r}(t)$  (2.12')  
(iii)  $\sum_{i} \xi_{i}(t) = 1$  and  $\sum_{i} \beta_{i}^{r}(t) = 1$ 

The fact that  $\xi_i(t) = \alpha_i(t) \beta_i(t) \ge 0$  and  $\sum_{i} \xi_i(t) = 1$  together with the above mentioned assumptions again allows us to use the information theory methodology to calculate the most

probable values  $\xi_i(t) = \hat{\xi}_i(t)$  (and thus  $\alpha_i(t) = \hat{\alpha}_i(t)$ ) by solving at any fixed time point the following minimization problem:

$$\begin{array}{l} \text{Min } \Sigma \xi_{1}^{r} \log \xi_{1}^{r} \\ \text{s.t.} \\ \xi_{1}^{r} P_{I}^{r} \leq \beta_{1}^{r}(t) P_{1}^{r}(\max) \end{array}$$

$$(2.14)$$

# 2.3 Composition of Trade Flows

Consider the trade flows originating from region r. As far as quality requirements and demand for attributes differ in two specific destination regions, s and k, we should in general expect  $x_I^{rs}$  to have a composition different from  $x_I^{rk}$ .<sup>2)</sup> Such a differentiation of flows is a natural extension of the heterogeneity of regional commodities  $x_I^{r*}$  as described in section 2.2.

<sup>2)</sup> The difference in composition means, e.g., that  $\beta_{i}^{rs} \neq \beta_{i}^{rk}$ ,  $i \in I$ .

Referring to formulae (2.2), (2.3) and (2.9) we shall assume that  $W_i^{rs}$  and  $V_i^{rs}$  are measured in such a way that flow prices,  $p_i^{rs}$ , can be calculated as  $p_i^{rs} = V_i^{rs}/x_i^{rs} = W_i^{rs}/x_i^{rs} - c_i^{rs}$ . Using the notation  $w_i^{rs} = W_i^{rs}/x_i^{rs}$  we shall examine the following four types of flow prices:

$$p_{i}^{rs} = w_{i}^{rs} - c_{i}^{rs}$$

$$p_{I}^{rs} = \sum_{i \in I} (w_{i}^{rs} - c_{i}^{rs}) x_{i}^{rs} / x_{I}^{rs} = w_{I}^{rs} - c_{I}^{rs}$$

$$p_{I}^{RS} = \sum_{r \in R} \sum_{s \in S} (w_{i}^{rs} - c_{i}^{rs}) x_{i}^{rs} / x_{i}^{RS} = w_{i}^{RS} - c_{i}^{RS}$$

$$p_{I}^{RS} = \sum_{r \in S} (w_{I}^{rs} - c_{I}^{rs}) x_{I}^{rs} / x_{I}^{RS} =$$

$$= \sum_{i \in I} (w_{i}^{RS} - c_{i}^{RS}) x_{i}^{RS} / x_{I}^{RS} = w_{I}^{RS} - c_{I}^{RS}$$

$$(2.15)$$

In Figure 3 we describe some of the possible disaggregation schemes that may be applied to an aggregate flow (I,R,S) characterized by  $p_{I}^{RS}$  and  $x_{I}^{RS}$ . The boxes in the figure with double borderlines indicate those levels of disaggregation on which our interest will be concentrated.

Let us illustrate the process of disaggregation by examining the following two cases:

(i) {
$$p_{I}^{RS}, x_{I}^{RS}$$
}  $\rightarrow$  { $p_{i}^{RS}, x_{i}^{RS}$ }, and  
(ii) { $p_{I}^{RS}, x_{I}^{RS}$ }  $\rightarrow$  { $p_{I}^{rs}, x_{I}^{rs}$ },

where  $\{p_{I}^{RS}, x_{I}^{RS}\}$  represents the outcome of a model projection for a future point in time.

Suppose that additional information which has not been utilized in the model projection is available. With regard to case (i), let this information be

o An a priori probability distribution  $\bar{\beta}_{i}^{RS}$ , indicating the expected likelihood that an arbitrary element in the flow  $x_{I}^{RS}$  belongs to commodity category i.



Figure 3. Illustration of disaggregation schemes.

- 0
- Estimates of transportation costs  $c_i^{RS}$ . Price scenarios  $p_i^{RS} = \alpha_i^{RS} p_I^{RS}$  for individual commodity flows  $x_i^{RS}$ . 0

Given this information, we may apply a minimum information principle to calculate the most probable values  $\beta_{i}^{RS} = x_{i}^{RS}/x_{I}^{RS}$ . This procedure may be summarized by the following Lagrange function

$$L = \sum_{i} \beta_{i}^{RS} \log \beta_{i}^{RS} / \beta_{i}^{RS} + \lambda \sum_{i} \beta_{i}^{RS} (1 - \alpha_{i}^{RS}) + \gamma [\sum_{i} \beta_{i}^{RS} c_{i}^{RS} - c_{I}^{RS}]$$

$$(2.16)$$

1) Compare Snickars and Weibull (1977).

This yields the solution

$$\beta_{i}^{RS} = \frac{\overline{\beta}_{i}^{RS} \exp \{\lambda \alpha_{i}^{RS} - \gamma c_{i}^{RS}\}}{\sum_{i} \overline{\beta}_{i}^{RS} \exp \{\lambda \alpha_{i}^{RS} - \gamma c_{i}^{RS}\}}$$
(2.17)

An analogous approach with regard to case (ii) yields

$$\beta_{I}^{rs} = \frac{\overline{\beta}_{I}^{rs} \exp \{\lambda \alpha_{I}^{rs} - \gamma c_{I}^{rs}\}}{\sum \sum \overline{\beta}_{I}^{rs} \exp \{\lambda \alpha_{I}^{rs} - \gamma c_{I}^{rs}\}}$$
(2.18)

Formulae (2.16) and (2.17) illustrate how alternative price scenarios affect the  $\beta_{I}^{rs}$ - and  $\beta_{i}^{RS}$ -distributions which are compatible with a given aggregate projection.

#### 3. HOMOGENEITY OF GROUPS IN AGGREGATION SCHEMES

An aggregation scheme consists of groups formed with regard to characteristics in three dimensions (i,r,s) where i  $\in$  II and r,s,  $\in$  R. As described in (2.1) a group is identified by a triple (I,R,S). A purposeful aggregation has to be constructed on the basis of a set of criteria, which is selected contingent on the modeling purpose. Here we shall confine the discussion to criteria which require that each group is homogenous in a problemspecific sense. Since we are considering projections into the future of aggregate flows  $x_I^{RS}$ , we want the homogeneity property of the elements  $x_i^{rs}$ , where (i,r,s) $\in$ (I,R,S), to be approximately time invariant.

#### 3.1 Direct and Generalized Commodity Homogeneity

In section 2.2 regional commodity groups were introduced. From (2.7) and (2.8) we can define

$$\mathbf{x}_{I}^{R*} = \sum_{S} \mathbf{x}_{I}^{RS} = \sum_{i \in I} \mathbf{x}_{i}^{R*}$$

$$\mathbf{x}_{I}^{*S} = \sum_{R} \mathbf{x}_{I}^{RS} = \sum_{i \in I} \mathbf{x}_{i}^{*S}$$
(3.1)

Homogeneity may be seen as a key to disaggregation of observed or known aggregates. For a regional supply group (I,R,\*) we may consider the following criterion:

Input homogeneity: this type of homogeneity is
satisfied if the input coefficient vector of
every good i∈I is proportional to the aggregate
input vector referring to I as a whole.
(3.2)

Condition (3.2) implies that cost conditions are similar for all i $\in$ I. It is easy to see that this condition only can be fulfilled if either (i) input coefficients are constant or (ii) production levels are proportional so that every ratio  $x_i^{R*}/x_T^{R*}$  is constant.

For a regional demand group (I,\*,S) we may consider two types of criteria. The first requires that all  $i \in I$  are "decomposable substitutes" in the sense specified below. The other requires that the goods  $i \in I$  are "decomposable complements".

> Observation 1: Suppose that all  $i \in I$  are perfect substitutes in the sense that there is only one market price  $w_I^{*S}$ . Then the goods are decomposable substitutes if there exist separable functions  $g_i^S$  such that  $x_i^{*S} = g_i^S(w_I^{*S})$  for each  $i \in I$ . If this is fulfilled all information about the disaggregation scheme  $\{w_i^{*S}, x_i^{*S}\}$  is contained in  $w_I^{*S}$ .

Observation 2: Suppose that all goods  $i \in I$  are perfect complements in that there exist coefficients such that  $\delta_i^{S} = x_i^{S}/x_I^{S}$ . These different goods are decomposable complements if there exist separable functions  $f_i^S$  such that  $w_i^{S} = f_i^S(x_i^{S})$ . If this is satisfied all information about  $\{x_i^{S}, w_i^{S}\}$  is contained in  $x_T^{S}$ .

Consider now the problem of aggregating commodities to a group I in such a way that the flow (R,S) is decomposable.

for criterion following the (2.12) yields decomposability: An extension of strong flow

$$\begin{aligned} x_{i}^{RS}(t) &= \beta_{i}^{RS} x_{I}^{SS}(t) \text{ for each } i \in I \\ p_{i}^{RS}(t) &= \alpha_{i}^{RS} p_{I}^{RS}(t) \text{ for each } i \in I \end{aligned} \tag{3.3}$$

 $\sim$ 

composition and the relative aggregate the individual the and pair  $(x_{I}^{RS}(t), p_{I}^{RS}(t)$  is a perfect proxy for flow terms  $\{p_{i}^{RS}(t), x_{i}^{RS}(t)\}$ . invariant, case the commodity time are simultaneously In this prices

Time invariance of coefficients guarantee decomposability ensures time invariant price ratios  $w_{i}^{RS}(t)/w_{I}^{RS}(t)$ ţ condition time invariance sufficient However, strong flow Observation 3:  $\alpha_{i}^{RS}$  is not a suf this latter

This observation follows directly from the definition  $w_{1}^{RS}(t) = \alpha_{1}^{RS} p_{I}^{RS}(t) + c_{1}^{RS}$  and  $w_{I}^{RS}(t) = p_{I}^{RS}(t) + \Sigma c_{1}^{RS} c_{1}^{RS}$  (t).

steps, starting and  $\rho_{ij}^{\alpha}$  where t t type described time series  $T = \{1, \dots, \overline{T}\}$  and wish and SCR with calculating correlation coefficients  $\rho^{g}_{ij}$  for a given complete spatial partitioning  $R^{C}R$ This could be done in hierarchical of invariances of the Suppose we have a investigate the degree (3.3). 'n

$$\rho_{ij}^{\beta} = \operatorname{Corr}_{t\in T} (x_i^{RS}(t), x_j^{RS}(t))$$

$$\nu \in T$$

$$\rho_{ij}^{\alpha} = \operatorname{Corr}_{\{R,S\}} (p_i^{RS}(t), p_j^{RS}(t))$$
(3.4)

specific (i) (3.4) into time-specific measures ( $\rho_{jj}^{\beta}(t), \rho_{jj}^{\alpha}(t)$ ) and (ii) flow measures ( $\rho_{jj}^{\beta}(R,S), \rho_{\alpha}(R,S)$ ) such that in We may split the correlation measures

$$\rho_{ij}^{\beta}(t) = \operatorname{Corr}_{\{R,S\}} (x_i^{RS}(t), x_j^{RS}(t)) \quad \text{for t given}$$

$$\rho_{ij}^{\beta (R,S)} = \operatorname{Corr}(x_i^{RS}(t), x_j^{RS}(t)) \text{ for } (R,S) \text{ given}$$

and analogously for  $\rho_{ij}^{\alpha}(t)$  and  $\rho_{ij}^{\alpha(R,S)}$ .

Observation 4: Let i, j \in I. If  $\alpha_i^{RS}$  and  $\alpha_j^{RS}$  are time invariant,  $\rho_{ij}^{\alpha(R,S)}$  will be close to one. If  $\alpha_i^{RS}$  and  $\alpha_j^{RS}$  are time invariant for each pair R,S in a complete spatial partitioning of  $\mathbb{R} \times \mathbb{R}$ ,  $\rho_{ij}^{\alpha}$  will be close to one. Observe finally that  $\rho_{ij}^{\alpha(R,S)}$  close to unity does not imply time invariance of  $\alpha_i^{RS}$  and  $\alpha_j^{RS}$ .

The first part of Observation 4 is self-evident. The second part is obvious if we note that  $p_i^{RS}(t)$  and  $p_j^{RS}(t)$  are perfectly correlated if  $p_i^{RS}(t) = a_i p_I^{RS}(t) + b$  and  $p_j^{RS}(t) = a_j p_I^{RS}(t) + b$ .<sup>1)</sup> However,  $\alpha_i^{RS}(t) = a_i + b/p_I^{RS}(t)$  will be time invariant only if  $p_I^{RS}(t)$  is constant.

Observation 5: The conclusions in Observation 4 are also true when  $\beta$  is substituted for  $\alpha$ .

Observation 6: Suppose that there is a time invariant relation between  $\beta_{i}^{RS}(t)$  and  $\beta_{j}^{RS}(t)$  such that  $\beta_{i}^{RS}(t) = b_{ij}\beta_{j}^{RS}(t)$  holds approximately for b fixed. Then  $\rho_{ij}^{\beta(R,S)}$  will be close to unity. Similarly,  $\rho_{ij}^{\alpha(R,S)}$  will be close to unity if  $\alpha_{i}^{RS}(t) = a_{ij}\alpha_{j}^{\alpha(t)}(t)$  holds approximately over time for  $a_{ij}$  fixed.

The conclusions in Observation 6 are obvious, since if  $x_i^{RS}(t) = \beta_i^{RS}(t)x_I^{RS}(t)$  for each  $i \in I$ , then  $x_j^{RS}(t)/x_i^{RS}(t) = 1/b_{ij}$  if  $i,j \in I$ .

<sup>1)</sup> This type of relation can only be valid when  $p_I^{RS}$  is constrained to satisfy  $p_I^{RS} \ge \pi$  for some  $\pi > 0$ .

Obviously we can also establish other forms of correlation and time-invariant properties. For example, if  $\beta_i^{RS}(t) = \overline{\beta}_i^{RS} f_i^{RS}(t)$ and  $\beta_i^{RS}(t) = \overline{\beta}_j^{RS} f_j^{RS}(t)$ , then we can examine the correlation between  $f_i^{RS}(t)$  and  $f_j^{RS}(t)$ , which is perfect if, for example,  $f_i^{RS}(t) = 1/f_j^{RS}(t)$ .

We may remark that the property of flow homogeneity has been transformed to find commodity groups within which there is a strong correlation between  $\beta_{i}^{RS}$ - and  $\alpha_{i}^{RS}$ -coefficients or time patterns of such coefficients, respectively.

#### 3.2 Procedures for Creating Commodity Groups

We shall describe a step-wise hierarchical procedure which starts from the fine level of commodity classification. In a first step each link (r,s) is examined.<sup>1)</sup> For each such link commodities are combined into commodity groups { $I^{1}(r,s),...,$  $I^{m}(r,s)$ } = I(r,s). Such a constellation of commodity groups is called a commodity classification.

In the subsequent step exporting regions r which have "similar" classifications are grouped together which generates new link pairs (R,s). Each such pair contains a set of classifications  $I(R,s) = \{I(r,s):r \in R\}$  which are similar but not necessarily identical. The final step of the exporter-based aggregation consists of finding for each pair (R,s) one single classification  $I^*(R,s)$  which can represent the set I(R,s) in a satisfactory way.

The three steps consist of finding (i) I(r,s), (ii) I(R,s), and (iii)  $I^*(R,s)$ . Before this a correlation matrix  $\{\rho_{ij}^{rs}\}$  is calculated for each given pair (r,s) and expressing the correlation between all possible pairs i,j $\in$ II. Referring to (3.6) we can calculate two versions of this matrix, namely  $\rho_{ij}^{\beta}(r,s)$ and  $\rho_{ij}^{\alpha}(r,s)$  2)

The objective is now to utilise the information in the correlation matrix for a given pair (r,s) in order to form groups containing commodities which strongly correlate with each

<sup>1)</sup> This procedure may also start with already spatially aggregated links (R,S) as in Secton 3.1.

<sup>2)</sup> Naturally it is also possible to calculate similar correlation coefficients with regard to transportation costs.

other. The following procedure will be outlined. We are searching for a threshold value  $\rho^*$  such that all pairs which simultaneously have pairwise values  $\rho_{ij} \ge \rho^*$  are included in the same group. Selecting  $\rho^* = 1$  means that the number of groups equals the total number of commodities. Reciprocally, selecting  $\rho^* = 0$  allows all commodities to be clustered together in one single group. In Figure 4 we have illustrated how the number of groups in a classification varies with the selection of a threshold value. The depicted curve makes a "jump" around the critical value  $\rho^{**}$ , i.e., in the neighbourhood of  $\rho^*$  curve  $N(\rho^*)$  has a significantly steeper slope than elsewhere;  $N(\rho^*)$  expresses the number of groups obtained when  $\rho^*$  is selected.



Figure 4. Number of groups in a classification as a function of  $\rho \textbf{*}.$ 

A desirable solution to the approach described in Figure 4 is obtained if we can find  $\rho^{**}$  which is close to one and which at the same time corresponds to a small number of groups in the classification. If the function  $N(\rho^*)$  increases without "jumps", for example, when  $N(\rho^*)$  has an almost constant first derivative, an exogenously determined compromise between the value of  $N(\rho^*)$  and  $\rho^*$  has to be utilized.

Suppose that  $\rho^{**}$  has been selected. Then the following approach may be applied:<sup>1)</sup>

<sup>1)</sup> Instead of this heuristic approach one may consider techniques like factor analysis.

- o Commodity i is compared with all  $j \neq i$ , and we select all j such that  $\rho_{ij}^{rs} \geq \rho^*$ ; these j's form a group I together with i.
- o Another commodity  $k \notin I$  is selected and the process is repeated.

With this procedure we have obtained one commodity classification I(r,s) for each pair (r,s),  $r,s \in \mathbb{R}$ .

The next step consists in bringing links (r,s) and (k,s)together in an export-related aggregate (R,s); the criterion for such spatial "clustering" of export regions (for given import region, s) is that the classifications I(k,s) and I(r,s)are similar for  $r,k\in R$ . In order to obtain a measure of similarity we first define

$$\theta_{ij}^{ks} = \begin{cases} 1 \text{ if commodity i and j belong to the} \\ \text{same group according to classification} \\ I(k,s) \\ 0 \text{ otherwise} \end{cases}$$
(3.7)

Hence,  $\theta_{ij}^{ks} = 1$  if  $i, j \in I$  and  $I \in I(k, s)$ . With this auxiliary variable we may calculate the distance between two classifications I(k, s) and I(r, s) as follows

$$d[I(k,s),I(r,s)] = \sum_{i,j\in II} \left| \theta_{ij}^{ks} - \theta_{ij}^{rs} \right|$$
(3.8)

Once these distances have been determined we may apply the same technique as was outlined for the correlation coefficients, and illustrated in Figure 4. The outcome will be (R,s)-specific commodity classifications as illustrated in Figure 5.

Our next task is to further compare the individual classifications  $I(\mathbf{r}, \mathbf{s}) \in I(\mathbf{R}, \mathbf{s})$ , where  $\mathbf{r} \in \mathbf{R}$ . On the basis of such a comparison we shall construct for each pair ( $\mathbf{R}, \mathbf{s}$ ) a universal classification  $I^{*}(\mathbf{R}, \mathbf{s})$  which will be used as a representative classification for ( $\mathbf{R}, \mathbf{s}$ ). This is accomplished by solving the following optimization problem:<sup>1</sup>

<sup>1)</sup> The algorithm for solving this problem is outlined in Section 4.1.



Figure 5. Illustration of I(R,s) and its elements I(r,s) for each importer s.

$$\min D ; \Sigma d[I^*(R,s), I(r,s)] = D$$
 (3.9)  
{I (R,s)}  $r \in R$ 

When this is solved we have obtained one classification for each pair (R,s).

At this stage we may identify a final aggregation problem which consists in forming groups of importers into importing regions. The outcome is pairs (R,S) and corresponding universal commodity classifications  $I^*(R,S)$ , such that individual importers  $s\in S$  are similar both with regard to exporter linkages and commodity classifications.

It should be observed that the hierarchical procedure makes it possible to collect structural information at each step of the procedure. Note also, that the aggregation may utilize more than one characteristic per link. Two such characteristics are the  $\alpha$ - and  $\beta$ -coefficients as defined in (2.12) and (3.3). Still one more characteristic is the transportation coefficients  $c_{i}^{rs}$ introduced in (2.2).

#### 3.3 Regional Aggregation with Given Commodity Groups

Consider the case of a given commodity i, or a given commodity group I. In the latter case we may, for example, perceive I as a group formed with all i and j such that each

$$\rho_{ij}^{\alpha} = Corr (p_{i}^{rs}(t), p_{j}^{rs}(t))$$
$$r, s \in \mathbb{R}$$
$$t \in T$$

has a high value, where  $\rho_{ij}^{\alpha}$  is the price correlation coefficient defined in (3.4).

In this subsection we restrict our analysis to one given commodity group. Therefore we can suppress or delete the commodity index. Our objective is to group together regions which simultaneously have similar export and import patterns as regards the given commodity. Each region may then be considered as a composite exporter and importer.

Trade patterns may be characterized as structural phenomena and reformulated as bivariate probability distributions for which information theory concepts are applicable (e.g., Andersson & Persson, 1982; Batten, 1982). For a given commodity group, this type of probability distribution is comprised of a matrix  $Q = \{q^{rs}\}$  where

$$q^{rs} = x^{rs} / \sum_{r,s \in \mathbb{R}} x^{rs}$$
(3.7)

and where  $x^{rs}$  denotes the flow between country r and s. The degree of dispersion or uncertainty in such a distribution may be calculated as  $\Sigma_{rs}q^{rs} \log q^{rs}$ . Transforming the matrix Q to another bivariate probability matrix Z then gives rise to the following information loss between Q and Z or change in the dispersion structure (see e.g., Snickars & Weibull, 1977): <sup>1)</sup>

$$I(Q:Z) = \sum_{r,s \in \mathbb{R}} q^{rs} \log (q^{rs}/z^{rs})$$
(3.8)

This type of measure will play a central role in the subsequent aggregation procedure. The aggregation of flows (r,s) to flows (R,S) gives us new probabilities

$$q^{RS} = \sum \sum q^{rs} , R, S \subset \mathbb{R}$$
(3.9)  
$$r \in R s \in S$$

<sup>1)</sup> This may be compared with the procedure in (2.16). Observe that I(Q:Z) may also be interpreted as the information gain between Z and Q.

Let  $\overline{Z} = {\{\overline{z}^{RS}\}}$  be an m x m-matrix and let the initial probability array Q be an n x n-matrix with m < n. Then we can find a permutation matrix n x m such that <sup>1)</sup>

$$\overline{z} = s^{\mathrm{T}} Q s \tag{3.10}$$

Every row of S has one unit element with all other elements being zeros. The non-zeros in each column show which elements are united into regional groups.

Consider now an auxiliary matrix Z with dimension n x n, which is "information equivalent" with  $\overline{Z}$ . We may define Z as a matrix obtained by

$$\begin{array}{l} \text{Min } \Sigma z^{\text{rs}} \log z^{\text{rs}} \\ z^{\text{rs}} \end{array}$$
 (3.11)

subject to summation constraints with regard to terms in Z which are associated with terms in  $\overline{Z}$ . In order to find an aggregate matrix  $\overline{Z}$  which has a "similar" structure to Q, we solve the following optimization problem (recalling the relation between Z and  $\overline{Z}$ ):

$$Min I(Q:Z) = Min \Sigma q^{rs} \log (q^{rs}/z^{rs})$$
(3.8)

where I(Q:Z) expresses the information loss between Q and Z caused by choosing Z instead of Q. As shown in Roy et al (1982) we can rewrite I(Q,Z) as follows:

$$I(Q:Z) = \sum_{r=1}^{n} \sum_{s=1}^{n} \log q^{rs} - \sum_{I=1}^{m} \sum_{J=1}^{m} \log (\overline{Z}_{IJ}/d(I)d(J))$$
  
r=1 s=1 I=1 J=1 IJ=1 IJ=1 IJ=1 (3.8')

where n and m are the number of rows (and columns) in the initial and aggregate matrices respectively, and where d(I) is the I'th element in the diagonal matrix D such that

$$D = S^{\mathrm{T}}S \tag{3.12}$$

1) S<sup>T</sup> denotes the transpose of S.

In this way I(Q:Z) is given as a direct function of the permutation matrix S. According to (3.10) this matrix is all we need in order to perform the aggregation.

Observation 7: Suppose that we have obtained aggregate probabilities (share coefficients)  $\bar{z}_{T}^{RS}$ . Then

- (i)  $\sum_{i} z_{i}^{rs} / \overline{z}_{I}^{RS}$  will represent  $\beta_{i}^{RS}$  as described in (2.17) and (3.3)
- (ii)  $\sum_{i \in I} z_{i}^{rs} / \overline{z}_{I}^{RS}$  will represent  $\beta_{I}^{rs}$  as described in (2.18).

## 4. AGGREGATION ALGORITHMS: Further Details

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In section 3.2 we described a procedure for aggregating commodities into classifications of commodity groups, aggregating exporter and/or importer regions, and identifying universal commodity classifications in the following steps

$$\{\rho_{ij}^{rs}\} \neq I(r,s) \neq I(R,s) \neq I^*(R,s)$$

If we like, this sequence can be continued further to establish a globally universal commodity classification  $I^*(*,*)$ , for which we can carry out a regional aggregation as described in section 3.3.

As an alternative, the procedure may start with a regional aggregation over a whole set of commodities, and thereafter construct I(R,S),  $I^*(R,S)$ , and eventually  $I^*(*,*)$ . Obviously, the search for universal classifications of different orders constitute a fundamental part of the aggregation.

# 4.1 An Algorithm for Finding Universal Classifications

For each pair (R,s) or (R,S) we are given a classification I(r,s) or I(r,S) for each  $r \in \mathbb{R}$ . The set of such classifications is I(R,s) or I(R,S). With regard to the pair (R,s) our objective is to find one classification  $I^*(R,s)$  that can universally

represent all classifications in I(R,s). This problem is posed in (3.9). Since the variables  $\theta_{ij}^{rs}$  in (3.9) are either unity or zero, the expression for D can be modified in the following way:

$$D = \sum_{\mathbf{r}} \sum_{i,j} |\theta_{ij}^{*s} - \theta_{ij}^{rs}| = \sum_{\mathbf{r}} \sum_{i,j} (\theta_{ij}^{*s} - \theta_{ij}^{rs})^{2} =$$

$$= \sum_{\mathbf{r}} \sum_{i,j} \theta_{ij}^{*s} - 2 \sum_{\mathbf{r}} \sum_{i,j} \theta_{ij}^{*s} \theta_{ij}^{rs} + \sum_{\mathbf{r}} \sum_{i,j} \theta_{ij}^{rs} = \qquad (4.1)$$

$$= \sum_{\mathbf{r}} \sum_{i,j} \theta_{ij}^{*s} - 2 \sum_{\mathbf{r}} \theta_{ij}^{*s} [\Sigma \theta_{ij}^{rs}] + \sum_{\mathbf{r}} \sum_{i,j} \theta_{ij}^{rs}$$

$$= \sum_{i,j} \theta_{ij}^{*s} - 2 \sum_{i,j} \theta_{ij}^{*s} [\Sigma \theta_{ij}^{rs}] + \sum_{i,j} \Sigma \theta_{ij}^{rs}$$

Let us introduce the variable  $a_{ij} = \Sigma_r \theta_{ij}^{rs}$  which is large for those pairs (i,j) which are included in the same group in many classifications and small in the opposite case. We may say that  $a_{ij}$  is a measure of "closeness" or "similarity" as regards (i,j). We may now rewrite (4.1) once more

$$D = N \Sigma \theta_{ij}^{*s} - 2 \Sigma \theta_{ij}^{*s} + \Sigma a_{ij}$$
(4.1')  
i,j i,j i,j

The last term can obviously be omitted from the optimization, since it occurs in (4.1') as a constant. Hence, the objective is to minimize

$$D = N \Sigma \theta_{ij}^{*s} - 2 \Sigma \theta_{ij}^{*s}$$
(4.2)

by selecting  $\{\theta_{ij}^{*s}\}$  in an optimal way.

Let the groups in classification  $I^* = I^*(\mathbf{R}, \mathbf{s})$  be indexed by 1,...,k,...,N so that  $I^* = \{I_1, \ldots, I_k, \ldots\}$ , and observe that  $\theta_{ij}^{*s} = 1$  for  $i, j \in I_k$  and that  $\theta_{ij}^{*s} = 0$  otherwise. Then the minimization problem in (4.2) may be formulated as

$$\max \sum_{\substack{k \text{ i, j} \in I_{k}}} \sum_{\substack{i \text{ i, j} \in I_{k}}} - (N/2) \sum N_{k}^{2}$$

$$(4.3)$$

where N<sub>k</sub> is the number of elements in  $I_k \in I^*$ . From formula (4.3)

we can see that two forces are operating: the objective function increases when the closeness component  $\Sigma\Sigmaa_{ij}$  grows; the function also increases when  $\Sigma N_k^2$  is reduced, i.e., when a uniform distribution is approached. N/2 defines the threshold of trade-off between the two conflicting goals. Having observed this we can replace the second term of (4.3) with a shift parameter  $0 \leq a \leq N^{-1}$  which expresses a selected threshold value. We can replace the optimization problem in (4.3) with the maximization of

$$U = \sum_{k i, j \in I_{k}} \sum_{k} (a_{ij} - a)$$
(4.4)

If a = 0 U reaches its maximum when all elements are placed in one group, and conversely if a = N the maximum obtains when every group contains just one element. Evidently, a can be used as shift or search parameter which is adjusted to give us the desired number of groups. A convenient start value may be a = N/2.

The algorithm consecutively unites individual elements or groups into gradually larger groups in such a way that the value of U increases for each change. This process is stopped when U does not show any further growth in value. The convergence of the process is guaranteed by the fact that U is bounded from above and at the same time monotonously nondecreasing.

4.2 An Algorithm for Finding the Permutation Matrix S

In (3.10) and (3.12) a permutation matrix is introduced. This transforms the initial trade pattern or disaggregated regional composition into an aggregate one. In (3.8) the solution was described as the outcome of the minimization of

$$I(Q,Z) = \Sigma q^{rs} \log q^{rs} - \Sigma \overline{Z}_{IJ} \log [\overline{Z}_{IJ}/d(I)d(J)]$$

The first term on the right hand side is not affected by the

<sup>1)</sup> N is the number of initial classifications.

selection of Z, and can therefore be deleted. Hence, the problem is reduced to the maximization of the second term. Recalling that d(I) and d(J) are elements in the diagonal matrix  $S^{T}S$ we can formulate this reduced problem as maximization of

$$\Sigma \mathbf{S}_{I}^{T} \mathbf{Q} \mathbf{S}_{J} \log \left[ \mathbf{S}^{T} \mathbf{Q} \mathbf{S}_{J} / \left( \mathbf{S}^{T} \mathbf{S} \right)_{I} \left( \mathbf{S}^{T} \mathbf{S} \right)_{J} \right]$$
(4.5)

where  $S_I$  and  $S_J$  are the I'th and J'th column vectors of the matrix S, respectively. In (4.5) the objective is entirely expressed in the sought matrix S and the known matrix Q. By specifying the desired number of regions, i.e., the dimensions of the matrix  $\overline{Z}$ , the function in (4.5) can be solved with some heuristic search procedure (see also Roy, Batten & Lesse, 1982).

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