

# **Convex Optimization by Radial Search**

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## **Working Paper**

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### Abstract

A convex nonsmooth optimization problem is replaced by a sequence of line search problems along recursively updated rays. Convergence of the method is proved and applications to linear inequalities, constraint aggregation and saddle point seeking indicated.

Key words: Nonsmooth optimization, subgradient methods, aggregation.

## Convex Optimization by Radial Search

Yuri M. Ermoliev Andrzej Ruszczyński

### 1 The method

The objective of this note is to present a new algorithmic concept for convex optimization problems of the form:

$$\min f(x), \quad x \in \mathbb{R}^n. \tag{1.1}$$

We assume that the function  $f: \mathbb{R}^n \mapsto \mathbb{R} \cup \{+\infty\}$  satisfies the following assumptions:

- (A1) f is convex, closed and co-finite, i.e.  $\sup_x \{\langle y, x \rangle f(x)\} < \infty$  for all  $y \in \mathbb{R}^n$ ;
- (A2)  $0 \in \text{int dom } f$ .

Consider the following method.

#### **ALGORITHM 1**

Step 0: Choose  $s^0 \in \mathbb{R}^n$  and  $\sigma \in (0,1)$ ; set k=0.

Step 1: Find  $x^k = -\mu_k s^k$  by minimizing f along the ray  $\{-\mu s^k : \mu \ge 0\}$ .

**Step 2:** Find a subgradient  $g^k \in \partial f(x^k)$  such that  $\left| \langle s^k, g^k \rangle \right| \leq \sigma |s^k|^2$  if  $x^k \neq 0$  and  $\langle s^k, g^k \rangle \leq \sigma |s^k|^2$  if  $x^k = 0$ .

Step 3: Set  $s^{k+1} = (1 - \tau_k)s^k + \tau_k g^k$ , increase k by one and go to Step 1.

Our method employs line search, as some of the bundle methods of [3, 4], but has a simple direction-generating rule, close to the subgradient averaging employed in some stochastic subgradient algorithms [1, 6]. Moreover, we do not increment  $x^k$  in successive directions, but we stay at one point (here 0) and we explore the space along selected rays. The method emerged from our recent work [2] on constraint aggregation schemes.

Throughout the paper we shall assume the following conditions on the stepsizes  $\{\tau_k\}$ .

- (A3)  $\tau_k \in [0,1];$
- (A4)  $\tau_k \to 0$ ;
- (A5)  $\sum_{k=0}^{\infty} \tau_k = \infty$ .

We shall base our analysis on the following lemma (see [2]).

**Lemma 1.1.** Let the sequences  $\{\beta_k\}$ ,  $\{\tau_k\}$ ,  $\{\delta_k\}$  and  $\{\gamma_k\}$  satisfy the inequality

$$0 \le \beta_{k+1} \le \beta_k - \tau_k \delta_k + \gamma_k. \tag{1.2}$$

If

- (i)  $\liminf \delta_k \geq 0$ ;
- (ii) for every subsequence  $\{k_i\} \subset \mathbb{N}$  one has  $[\liminf \beta_{k_i} > 0] \Rightarrow [\liminf \delta_{k_i} > 0]$ ;
- (iii)  $\tau_k \geq 0$ ,  $\lim \tau_k = 0$ ,  $\sum_{k=0}^{\infty} \tau_k = \infty$ ;
- (iv)  $\lim \gamma_k/\tau_k = 0$ ,

then  $\lim_{k\to\infty} \beta_k = 0$ .

**Proof.** Suppose that  $\liminf \delta_k = \delta > 0$ . Then (1.2) for large k yields  $\beta_{k+1} \leq \beta_k - \tau_k \delta/2 + \gamma_k \leq \beta_k - \tau_k \delta/4$ . This contradicts (iii). Therefore  $\liminf \delta_k = 0$ . By (ii) there is a subsequence  $\{k_i\}$  such that  $\beta_{k_i} \to 0$ . Suppose that there is another subsequence  $\{s_j\}$  such that  $\beta_{s_j} \geq \beta > 0$  for  $j = 0, 1, 2, \ldots$  With no loss of generality we may assume that  $k_1 < s_1 < k_2 < s_2 \ldots$  By (i), (iii) and (iv), for all sufficiently large j there must exist indices  $r_j \in [k_j, s_j]$  such that  $\beta_{r_j} > \beta/2$  and  $\beta_{r_{j+1}} > \beta_{r_j}$ . But then, by (ii),  $\liminf \delta_{r_j} = \delta > 0$  and we obtain a contradiction with (1.2) for large j.  $\square$ 

**Lemma 1.2.** There exists a constant C such that for all k one has  $|g^k| \leq C(1 + |s^k|)$ .

**Proof.** Denote  $f_{\min} = \min f(x)$ . By (A2),  $f_{\min} > -\infty$ . For every  $\epsilon > 0$  we have

$$f\left(\frac{\epsilon g^k}{|g^k|}\right) \ge f(x^k) + \left\langle \frac{\epsilon g^k}{|g^k|} + x^k, g^k \right\rangle$$
$$\ge f_{\min} + \epsilon |g^k| - \mu_k \langle s^k, g^k \rangle.$$

Using the conditions of Step 2 we obtain

$$f\left(\frac{\epsilon g^k}{|g^k|}\right) \ge f_{\min} + \epsilon |g^k| - \sigma \mu_k |s^k|^2.$$

By (A2), the set  $X_0 = \{x \in \mathbb{R}^n : f(x) \leq f(0)\}$  has a finite diameter d. Therefore  $\mu_k |s^k| \leq d$ . Moreover, f is finite around 0, so for some small but fixed  $\epsilon > 0$  and some  $C_1$ ,  $f(\epsilon g^k / |g^k|) \leq C_1$  for all k. The last inequality then implies that

$$\epsilon |g^k| \le C_1 - f_{\min} + \sigma d|s^k|,$$

which yields the required result. \(\sigma\)

Lemma 1.3.  $\lim_{k \to \infty} s^k = 0$ .

**Proof.** By the conditions of Step 2,

$$|s^{k+1}|^2 = (1 - \tau_k)^2 |s^k|^2 + 2\tau_k (1 - \tau_k) \langle s^k, g^k \rangle + \tau_k^2 |g^k|^2$$
  

$$\leq (1 - 2(1 - \sigma)\tau_k + \tau_k^2) |s^k|^2 + \tau_k^2 |g^k|^2.$$
(1.3)

By Lemma 1.1,

$$|g^k|^2 \le C^2(1+|s^k|)^2 \le 2C^2(1+|s^k|^2).$$

Therefore,

$$|s^{k+1}|^2 \le (1 - 2(1 - \sigma)\tau_k + (2C^2 + 1)\tau_k^2)|s^k|^2 + 2C^2\tau_k^2$$

By (A4), for all sufficiently large k one has  $\tau_k \leq (1-\sigma)/(2C^2+1)$ , so

$$|s^{k+1}|^2 \le (1 - (1 - \sigma)\tau_k)|s^k|^2 + 2C^2\tau_k^2.$$

The required result follows now from Lemma 1.1. □

**Theorem 1.4.** Assume (A1)-(A5). Then for the sequence  $\{x^k\}$  generated by Algorithm 1 one has

$$\liminf f(x^k) = \min_{x \in \mathbb{R}^n} f(x).$$

**Proof.** Consider the conjugate function  $f^*(\cdot) = \max_x \{\langle x, \cdot \rangle - f(x)\}$  (see, e.g., [3, 5]). It is convex and (by assumption) finite everywhere. From the convexity of  $f^*$  we get

$$f^*(s^{k+1}) \le (1 - \tau_k) f^*(s^k) + \tau_k f^*(g^k).$$

From Fenchel's equality (see, e.g. [5, Thm. 23.5]) and conditions of Step 2 we obtain

$$f^*(g^k) = -f(x^k) + \langle x^k, g^k \rangle$$

$$= -f(x^k) - \mu_k \langle s^k, g^k \rangle$$

$$\leq -f(x^k) + \mu_k \sigma |s^k|^2$$

$$\leq -f(x^k) + \sigma d|s^k|,$$

where d is the upper bound on  $|x^k| = \mu_k |s^k|$ . Combining the last two inequalities we obtain

$$f^*(s^{k+1}) \le f^*(s^k) - \tau_k(f^*(s^k) + f(x^k) - \sigma d|s^k|). \tag{1.4}$$

By the continuity of  $f^*$ ,  $f^*(s^k) \to f^*(0) = -f_{\min}$ . Suppose that  $f(x^k) \ge f_{\min} + \epsilon$  for all k, where  $\epsilon > 0$ . Then (1.4), Lemma 1.3 and (A5) imply that  $f^*(s^k) \to -\infty$ , a contradiction. Therefore  $\liminf f(x^k) = f_{\min}$ .  $\square$ 

A stronger result can be obtained for the sequence of averages.

**Theorem 1.5.** Let the assumptions of Theorem 1.4 be satisfied. Then for the sequence of averages

$$\bar{x}^{k+1} = (1 - \tau_k)\bar{x}^k + \tau_k x^k, \quad k = 0, 1, 2, \dots,$$

where  $\{x^k\}$  is generated by Algorithm 1, one has

$$\lim_{k \to \infty} f(\bar{x}^k) = \min_{x \in \mathbb{R}^n} f(x).$$

**Proof.** From the convexity of f and  $f^*$  we obtain

$$f(\bar{x}^{k+1}) \le (1 - \tau_k)f(\bar{x}^k) + \tau_k f(x^k),$$

$$f^*(s^{k+1}) \le (1 - \tau_k) f^*(s^k) + \tau_k f^*(g^k).$$

Adding both sides yields

$$f(\bar{x}^{k+1}) + f^*(s^{k+1}) \le (1 - \tau_k)(f(\bar{x}^k) + f^*(s^k)) + \tau_k \langle x^k, g^k \rangle.$$

because  $f(x^k) + f^*(g^k) = \langle x^k, g^k \rangle$  [5, Thm. 23.5]. By the conditions of Step 2,  $\langle x^k, g^k \rangle \le \mu_k \sigma |s^k|^2 \le d|s^k|$ , where d is the upper bound on  $|x^k|$ . Therefore,

$$\max(0, f(\bar{x}^{k+1}) + f^*(s^{k+1})) \le (1 - \tau_k) \max(0, f(\bar{x}^k) + f^*(s^k)) + \tau_k d|s^k|.$$

Since  $|s^k| \to 0$  by Lemma 1.3, using Lemma 1.1 we conclude that

$$\lim_{k \to \infty} \max(0, f(\bar{x}^k) + f^*(s^k)) = 0. \tag{1.5}$$

With  $f^*(s^k) \to f^*(0) = -f_{\min}$ , the required result follows from (1.5).  $\Box$ 

## 2 Explicit non-negativity constraints

The concept introduced in section 1 applies, of course, to constrained problems, because we allow  $+\infty$  as the value of f. For example, simple inequalities  $x \ge 0$  can be dealt with by moving the center 0 to some  $\tilde{x} > 0$ . It is, however, more convenient to treat them explicitly.

Consider the problem

$$\min_{x>0} f(x)$$

under the same assumptions as before. Then we can still apply the method described in section 1, with the following modifications.

#### **ALGORITHM 2**

**Step 0:** Choose  $s^0 \in \mathbb{R}^n$  and  $\sigma \in (0,1)$ ; set k=0.

Step 1: Find  $x^k = \mu_k d^k$  by minimizing f along the ray  $\{\mu d^k : \mu \geq 0\}$ , where  $d^k$  is the projection of  $-s^k$  onto the positive orthant:  $d_j^k = \max(0, -s_j^k), j = 1, \ldots, n$ .

**Step 2:** Find a subgradient  $g^k \in \partial f(x^k)$  such that  $\left| \langle d^k, g^k \rangle \right| \leq \sigma |d^k|^2$  if  $x^k \neq 0$  and  $\langle d^k, g^k \rangle \geq -\sigma |d^k|^2$  if  $x^k = 0$ .

Step 3: Set  $s^{k+1} = (1 - \tau_k)s^k + \tau_k g^k$ , with  $\tau_k \in [0, 1]$ , increase k by one and go to Step 1.

The convergence properties remain unchanged.

**Theorem 2.1.** Let the assumptions of Theorem 1.4 be satisfied. Then for the sequence  $\{x^k\}$  generated by Algorithm 2 one has

$$\liminf f(x^k) = \min_{x \ge 0} f(x).$$

**Proof.** We shall derive a counterpart of the key inequality (1.3). From the definition of  $d^k$  one obtains

$$-s^{k+1} \le (1 - \tau_k)d^k - \tau_k g^k.$$

In the above vector inequality, for the components j such that  $-s_j^{k+1} > 0$  the absolute value of the right hand side is not less than  $|s_i^{k+1}|$ , so

$$|d^{k+1}|^2 \le |(1-\tau_k)d^k - \tau_k g^k|^2$$

$$= (1-\tau_k)^2 |d^k|^2 + 2\tau_k (1-\tau_k) \langle d^k, g^k \rangle + \tau_k^2 |g^k|^2$$

$$\le (1-2(1-\sigma)\tau_k + \tau_k^2) |d^k|^2 + \tau_k^2 |g^k|^2,$$

where in the last inequality we used the conditions of Step 2. Proceeding exactly as in the proofs of Lemmas 1.2 and 1.3, we conclude that  $d^k \to 0$  and  $\{g^k\}$  is bounded. Then the sequence of averages  $\{s^k\}$  is bounded, too. Let  $\bar{s}$  be any accumulation point of  $\{s^k\}$ . Since  $d^k \to 0$ , one must have  $\bar{s} \geq 0$ . By the continuity of  $f^*$ , for the corresponding subsequence we get

$$\begin{split} f^*(s^k) &\to f^*(\bar{s}) = \max_x \{ \langle \bar{s}, x \rangle - f(x) \} \\ &\geq \max_{x \geq 0} \{ \langle \bar{s}, x \rangle - f(x) \} \geq -f_{\min}, \end{split}$$

where  $f_{\min} = \min_{x \geq 0} f(x)$ . Consequently,

$$\liminf f^*(s^k) \ge -f_{\min}. \tag{2.1}$$

This combined with inequality (1.4), in the same manner as in Theorem 1.4, yields the required result.  $\Box$ 

We also have an analog of Theorem 1.5.

**Theorem 2.2.** Let the assumptions of Theorem 1.4 be satisfied. Then for the sequence of averages

$$\bar{x}^{k+1} = (1 - \tau_k)\bar{x}^k + \tau_k x^k, \quad k = 0, 1, 2, \dots,$$

where  $\{x^k\}$  is generated by Algorithm 2, one has

$$\lim_{k \to \infty} f(\bar{x}^k) = \min_{x > 0} f(x).$$

**Proof.** Proceeding similarly to the proof of Theorem 1.5 we obtain relation (1.5), which implies

$$\limsup \left( f(\bar{x}^k) + f^*(s^k) \right) \le 0. \tag{2.2}$$

On the other hand,  $f(\bar{x}^k) \geq f_{\min}$ , so we must have  $\limsup f^*(s^k) \leq -f_{\min}$ . This combined with (2.1) yields

$$\lim_{k \to \infty} f^*(s^k) = -f_{\min}.$$

Our assertion follows now from (2.2).  $\square$ 

### 3 Applications

Let us discuss some potential applications of the ideas introduced in this paper.

Linear inequalities

Consider the system of linear inequalities

$$\sum_{j=1}^{n} a_{ij} x_j \le b_i, \ i = 1, \dots, m, \tag{3.1}$$

and the associated optimization problem

$$\min_{x} \left[ f(x) = \max_{1 \le i \le m} \left( \sum_{j=1}^{n} a_{ij} x_j - b_i \right) \right].$$

The subproblem solved at Step 1 takes on the form

$$\min_{\mu \geq 0} \max_{1 \leq i \leq m} \left( -\mu \sum_{j=1}^{n} a_{ij} s_j^k - b_i \right).$$

Define the sets

$$J_k^+ = \{j : \sum_{j=1}^n a_{ij} s_j^k > 0\},$$

$$J_k^- = \{j: \sum_{i=1}^n a_{ij} s_j^k \le 0\}.$$

If  $J_k^- = \emptyset$  then  $\sum_{j=1}^n a_{ij} s_j^k > 0$  for all i and one can find  $\bar{\mu} \ge 0$  such that  $-\bar{\mu} s^k$  solves (3.1). It remains to consider the case when  $J_k^- \ne \emptyset$  for all k.

If  $\mu_k > 0$  there must exist  $r \in J_k^-$  and  $t \in J_k^+$  such that

$$f(-\mu_k s^k) = -\mu_k \sum_{j=1}^n a_{rj} s_j^k - b_r = -\mu_k \sum_{j=1}^n a_{tj} s_j^k - b_t.$$

Denote  $a_r = (a_{r1}, \ldots, a_{rn}), a_t = (a_{t1}, \ldots, a_{tn})$  and define

$$\lambda_k = \frac{\langle a_t, s^k \rangle}{\langle a_t - a_r, s^k \rangle}.$$

Since  $a_r \in \partial f(x^k)$ ,  $a_t \in \partial f(x^k)$  and  $\lambda_k \in [0,1]$ ,

$$g^k = \lambda_k a_r + (1 - \lambda_k) a_t$$

is a subgradient of f at  $x^k$ . By the definition of  $\lambda_k$ ,  $\langle s^k, g^k \rangle = 0$ , i.e.  $g^k$  satisfies the conditions of Step 2 with  $\sigma = 0$ .

If  $\mu_k = 0$ , then there must exist  $r \in J_k^-$  such that  $b_r \leq b_i$ ,  $i = 1, \ldots, m$ . Taking  $g^k = a_r$  we have  $\langle g^k, s^k \rangle \leq 0$  by the definition of  $J_k^-$ .

Constraint aggregation

Consider the convex optimization problem

$$\min h(y) \tag{3.2}$$

$$Ay = b, (3.3)$$

$$y \in Y, \tag{3.4}$$

where  $h: \mathbb{R}^m \to \mathbb{R}$  is convex,  $Y \subset \mathbb{R}^m$  is convex and compact, A is an  $n \times m$  matrix,  $b \in \mathbb{R}^n$ . Its dual has the form

$$\max f(x), x \in \mathbb{R}^n,$$

where x is the vector of Lagrange multipliers and  $f: \mathbb{R}^n \mapsto \mathbb{R}$  is the dual function defined as follows:

$$f(x) = \min_{y \in Y} \left\{ h(y) + \langle x, Ay - b \rangle \right\}.$$

Clearly, -f is convex and co-finite. Let us apply Algorithm 1 to the dual problem (with obvious modifications reflecting the change from minimization to maximization). Step 1 takes on the form

$$\max_{\mu>0} \min_{y \in Y} \left\{ h(y) + \mu \langle s^k, Ay - b \rangle \right\},\,$$

which, under appropriate constraint qualification, is equivalent to the following optimization problem

$$\min h(y) \tag{3.5}$$

$$\langle s^k, Ay - b \rangle \le 0, \tag{3.6}$$

$$y \in Y. \tag{3.7}$$

The subgradient  $g^k$  satisfying the conditions of Step 2 is given by

$$g^k = Ay^k - b, (3.8)$$

where  $y^k$  is the solution of (3.5)-(3.7). Finally, the subgradient averaging rule of Step 3 can be written as

$$z^{k+1} = (1 - \tau_k)z^k + \tau_k y^k, (3.9)$$

$$s^{k+1} = Az^{k+1} - b. (3.10)$$

The algorithm (3.5)-(3.10) can be regarded as an iterative constraint aggregation procedure for solving (3.2)-(3.4): it replaces the constraints (3.3) by a single surrogate inequality (3.6). This idea has been analysed in [2].

If the original problem, instead of (3.3), has inequality constraints

$$Ay \leq b$$
,

the dual problem has non-negativity constraints on x, so Algorithm 2 applies. The only modification with respect to (3.5)-(3.10) is that (3.10) is replaced by the projection:

$$s^{k+1} = \left(Az^{k+1} - b\right)_+,$$

where  $(v_+)_j = \max(0, v_j)$ , j = 1, ..., n. In a similar way we can treat convex inequalities (see [2] for the details missing here, such as the constraint qualification condition, various modifications and extension, analysis of the rate of convergence, etc).

#### Saddle point seeking

The previous example can be in a straighforward manner generalized to the saddle point problem. Let  $L: \mathbb{R}^n \times Y \mapsto R$  be a convex-concave function. Assuming that L is strictly concave in its second argument, we can find a saddle point  $(\hat{x}, \hat{y})$  of L in the following way. First, we solve the problem

$$\min_{x \in \mathbb{R}^n} \left[ f(x) = \sup_{y \in Y} L(x, y) \right]$$
 (3.11)

to get  $\hat{x}$  and then we define  $\hat{y}$  as the maximizer of  $L(\hat{x},\cdot)$  over Y. It turns out that Step 1 of Algorithm 1 applied to (3.11) takes on the form:

$$\min_{\mu \geq 0} \sup_{y \in Y} L(-\mu s^k, y).$$

By defining the function  $\Lambda_k(\mu, y) = L(-\mu s^k, y)$  we can equivalently formulate Step 1 as follows: find a saddle point  $(\mu_k, y^k)$  of  $\Lambda_k$  on  $\mathbb{R}_+ \times Y$ . Moreover, if L is continuously differentiable with respect to the first argument, then  $g^k = \nabla_x L(-\mu_k s^k, y^k)$  satisfies the conditions of Step 2 with  $\sigma = 0$ .

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